

Static analysis for data race detection

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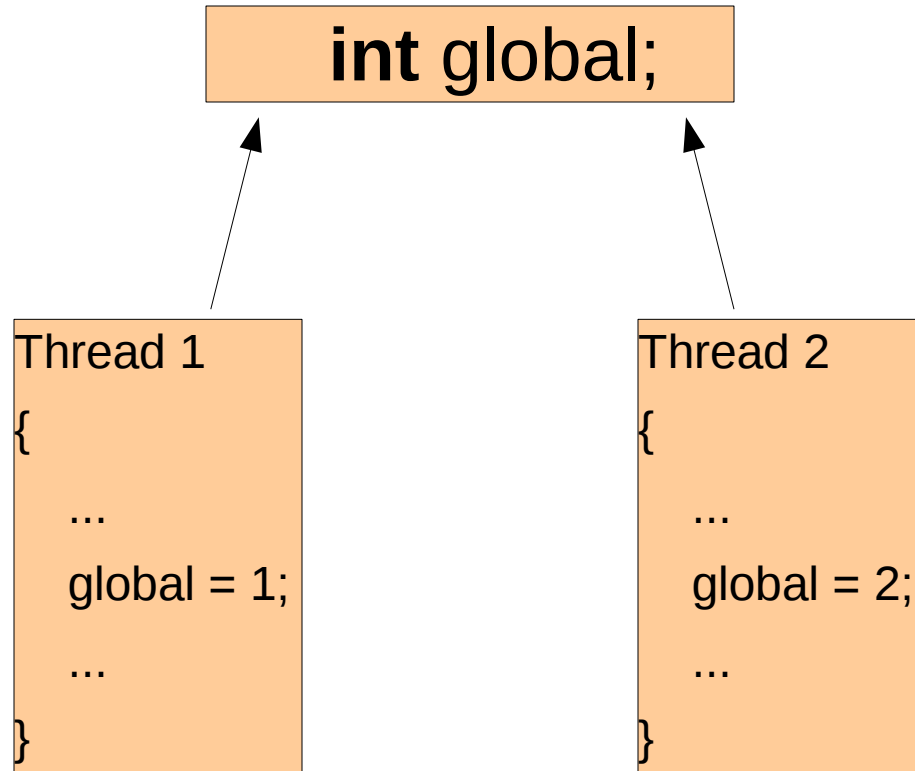
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Data Race Condition



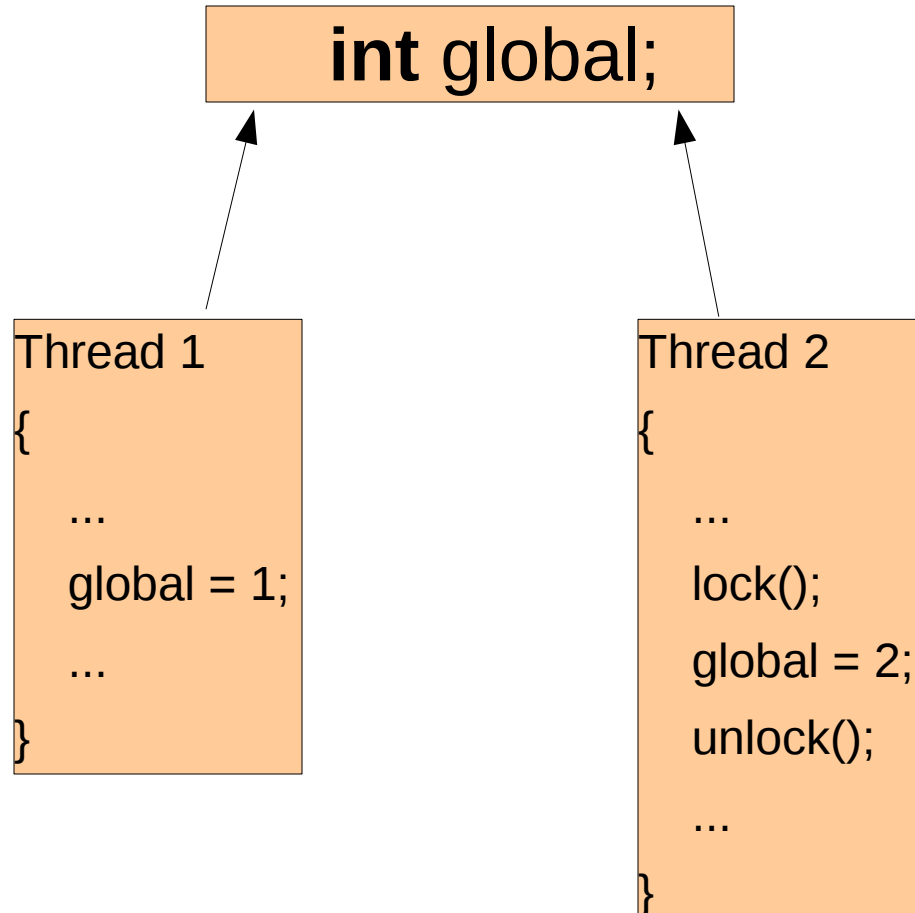
Notions

- ***Variable*** – local, global, structure field
- ***Lock*** – synchronization primitive
- ***Usage of variable*** – write or read data access
- ***Shared data*** – available from several threads data
- ***Unsafe*** – potential error (race)

Data Race Condition

- *Potential data race condition* is a situation, when usage of the same variable occurs with **different disjoint sets** of locks, one usage is write access.

Data Race Condition



Example

```
int func1() {  
    struct A *s;  
    int local;  
    s->f = ...  
    lock();  
    global = ...  
    local = ...  
    unlock();  
}
```

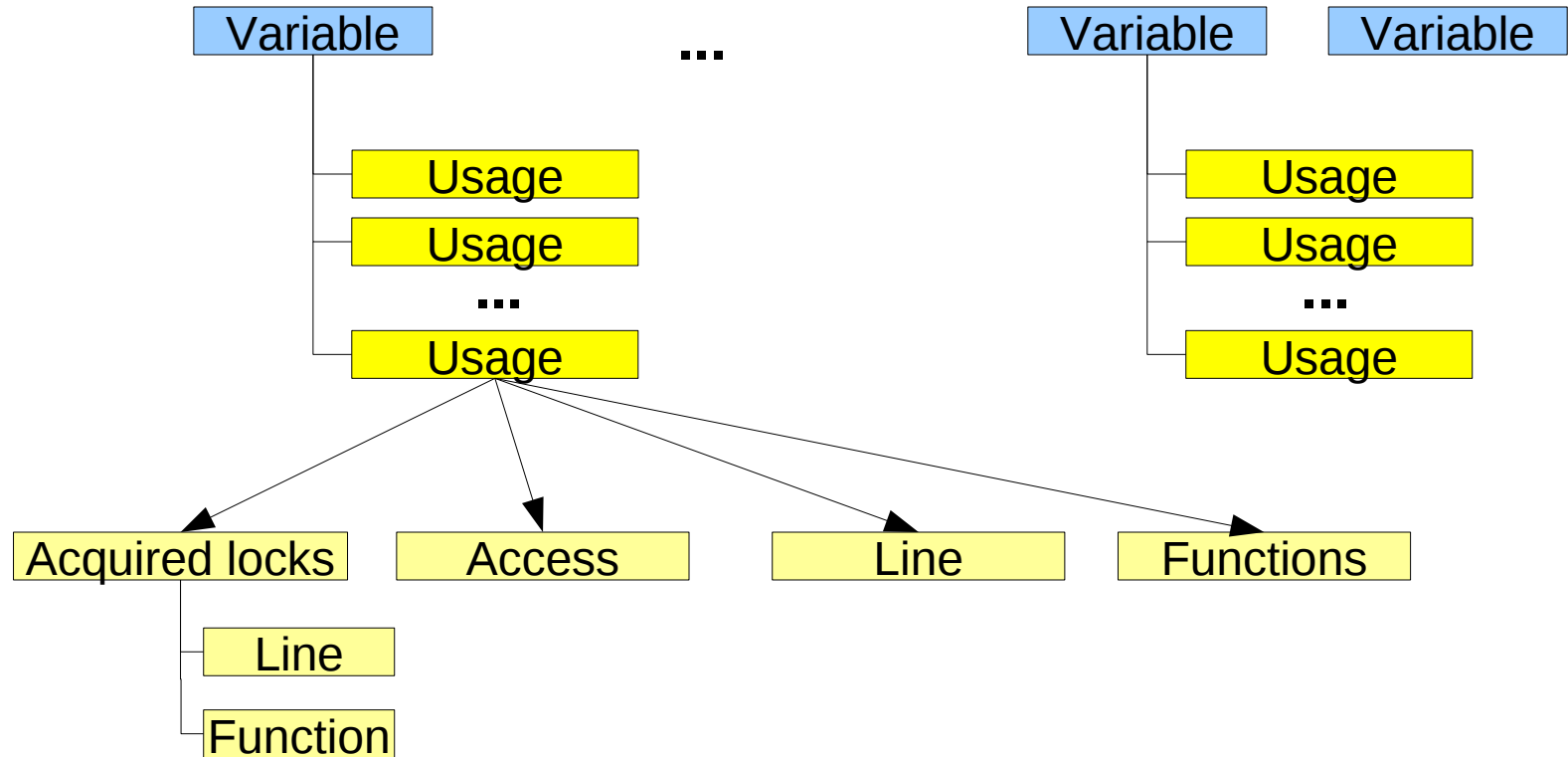
```
int func2() {  
    int local;  
    ...  
    local = global;  
}
```

Acquiring a lock
Releasing a lock
Write access to a variable
Read access to a variable

Potential race condition



Necessary information



Analysis of shared data

```
struct my_struct {  
    int* b;  
} A;  
int func() {  
    int *a;  
    a = malloc(); ← a -> local data  
    *a = 1; ← Write access to local data – not a race  
    A->b = a; ← a -> shared data  
    *a = 1; ← Write access to shared data – subject for data race  
    a = g(); ← a -> unknown data  
}
```

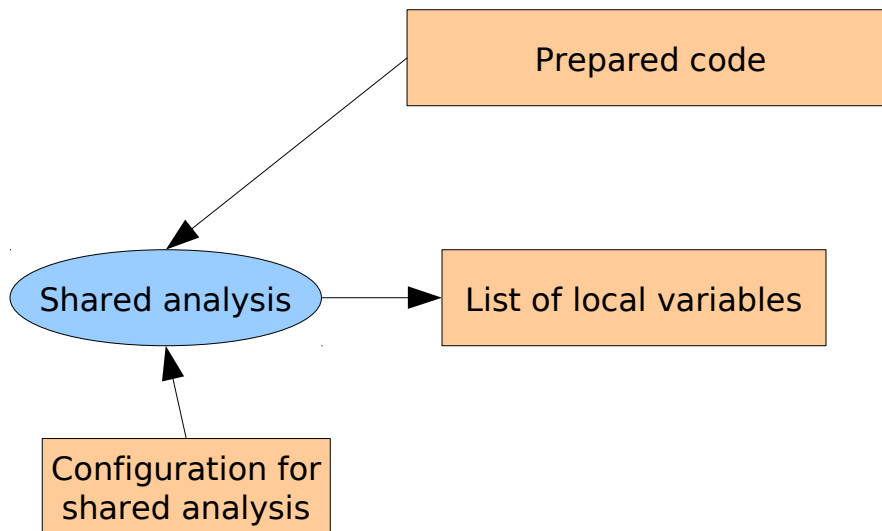

Stage of analysis: Instrumentation

- Replacement macroexpansions by calls of model functions
- Insert entry point (main function)

Prepared code

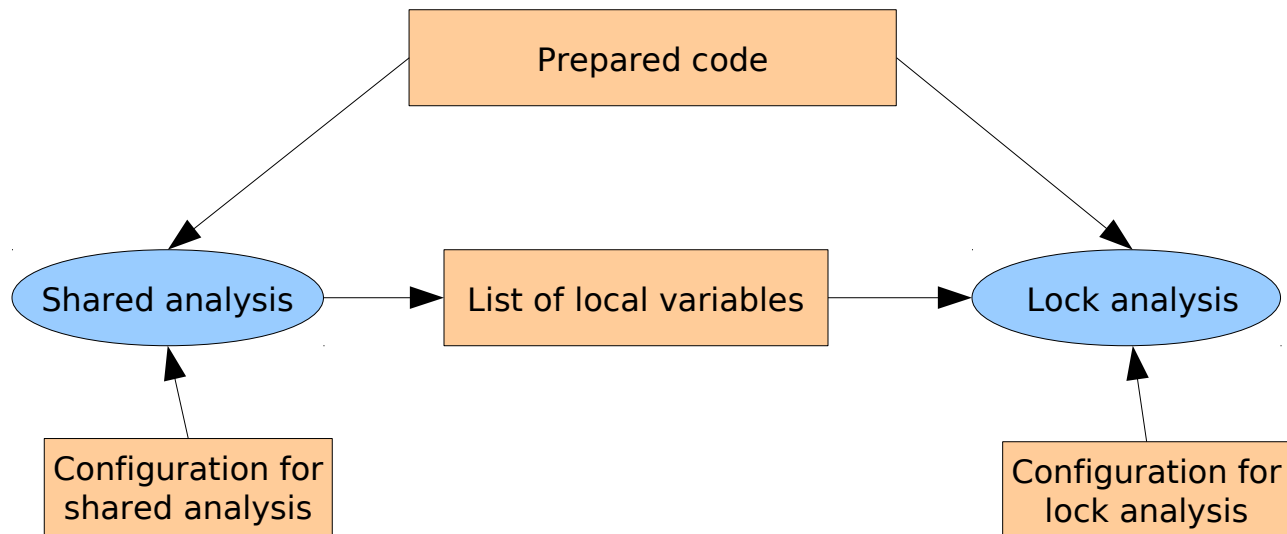
Stage of analysis: Analysis of shared data

- Obtain information about local variables for every point in program



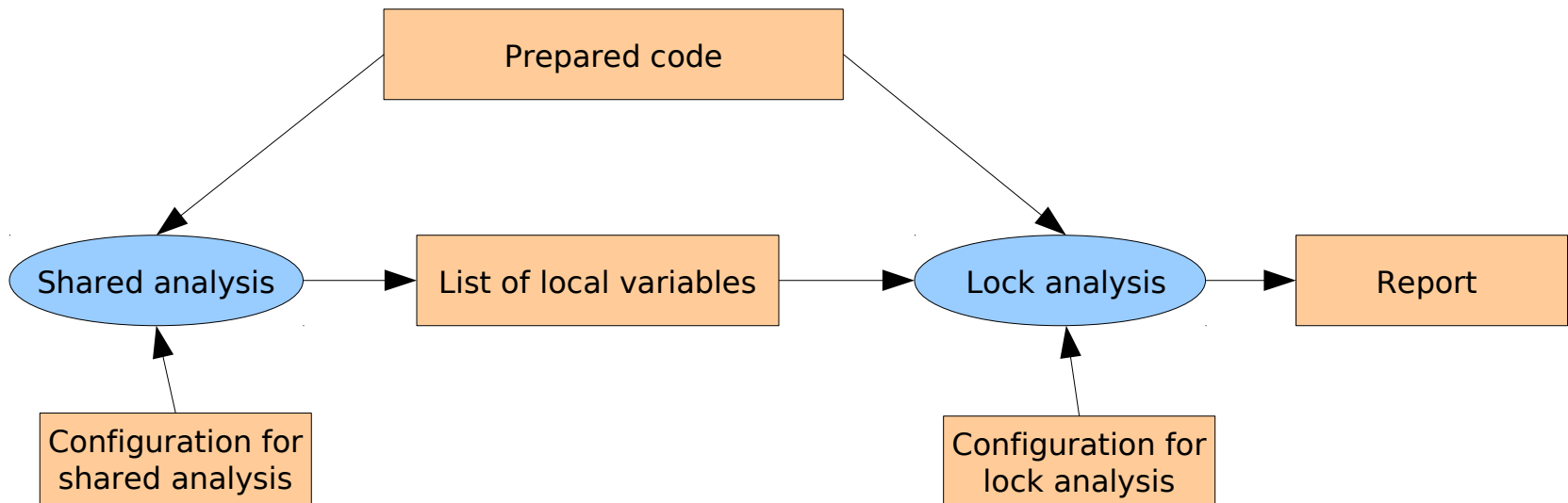
Stage of analysis: Analysis of synchronization primitives

Save information about acquired locks for every access to memory



Stage of analysis: Result

- List of unsafe cases
- Detail information about every unsafe



Report

Statistics	General	Unsafe
Global variables:	591	24
Simple:	411	20
Pointer:	180	4
Local variables:	1296	50
Simple:	0	0
Pointer:	1296	50
Structure fields:	1324	300
Simple:	1038	287
Pointer:	286	13
Total variables:	3211	374

Finded locks:

1. `global_lock lock()`
2. `global_lock pthread_mutex_lock(m->_mutex)(0)`
3. `global_lock pthread_mutex_lock(mutex)(0)`

List of unsafes:

Visualization of Unsafes

int *a

The image displays a code visualization tool with two main panels: "Error trace" and "Source code".

Error trace: This panel shows a call stack with the following details:

- Buttons: Function bodies, Blocks, Others...
- Code:

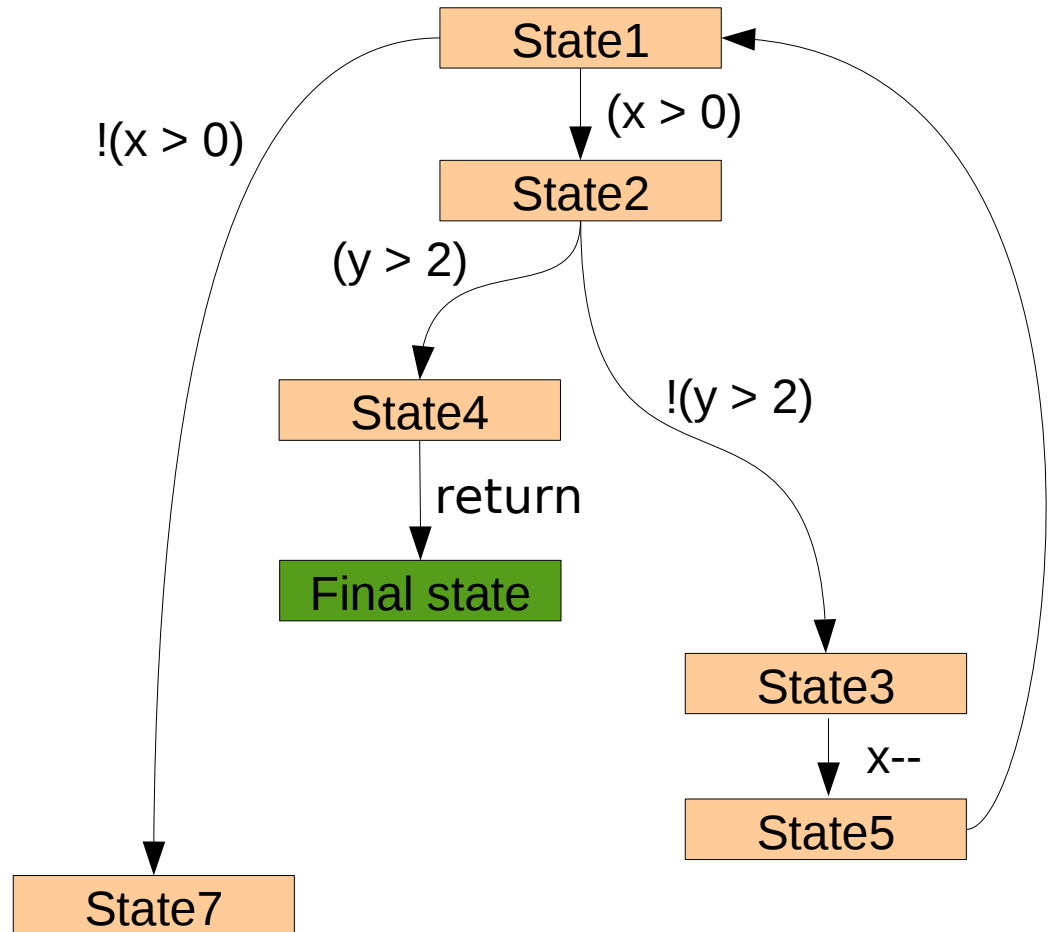
```
/*Number of usages:4*/  
/*Two examples:*/  
/*  
/*global_lock pthread_mutex_lock(m)(0)*/  
_main()  
{  
16 pthread_mutex_lock() { /* Function call is skipped d  
17 _f()  
{  
9 _g()  
{  
5 *a = ...;  
}  
}  
}  
/*Without locks*/  
_main()  
{  
17 _f()  
{  
9 _g()  
{  
5 ... = *a;  
}  
}  
}
```

Source code: This panel shows the source code for "Test.c":

```
1 #line 2 "/home/alpha/git/cpachecker/test/Test.c"  
2 int global;  
3  
4 int g(int *a) {  
5     (*a)++;  
6 }  
7  
8 int f(int *a) {  
9     int r = g(a);  
10    return r;  
11 }  
12  
13 int main() {  
14     mutex m;  
15     f(&global);  
16     pthread_mutex_lock(m);  
17     f(&global);  
18     pthread_mutex_unlock(m);  
19 }
```

Main idea of static analysis

```
while (x > 0)
{
    if (y > 2) {
        return;
    }
    x--;
}
```



Configurable Program Analysis (CPA)

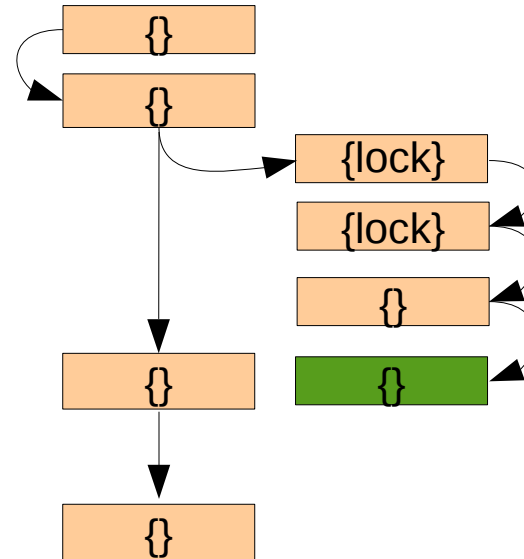
- Abstract domain – set of abstract states
- Transfer relation – for every state e it defines next state e'
- Merge operator – combines information about two abstract states
- Stop operator – checks, if current state e is covered by set of others

Example: Lock statistics CPA

- Abstract domain – set of acquired locks
- Transfer relation – changes state, when function of acquiring or releasing lock is called
- Merge operator – states never combines
- Stop operator – state is covered, only if given set contains equal state

Example: LockStatisticsCPA

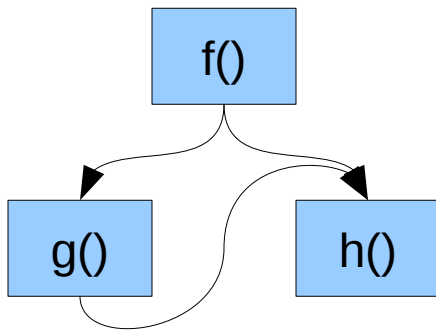
```
int global;  
int func() {  
    if (global) {  
        lock();  
        global++;  
        unlock();  
    }  
}
```



Main problems

- Resources: time and memory
- Analysis of infeasible paths

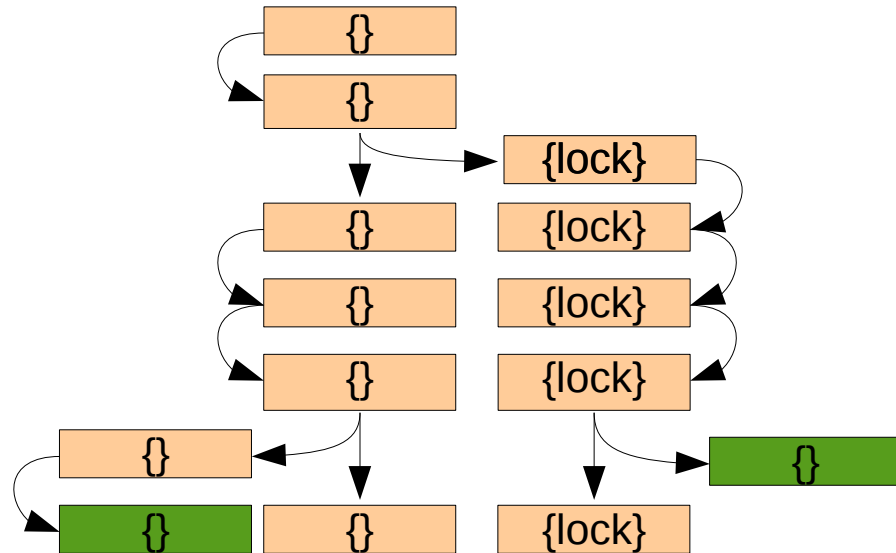
Abstract block memorization (ABM)



- Store the results of previous analysis
- Configurable blocks
- Greatly accelerate the analysis

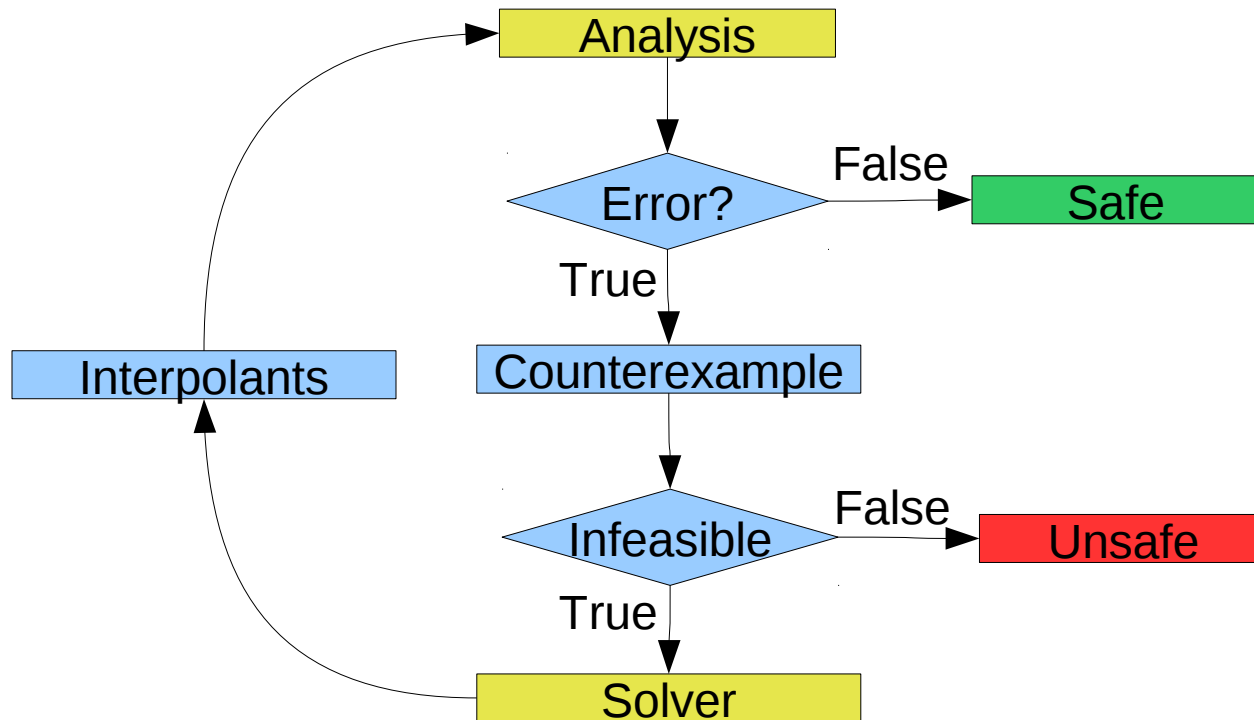
Example of infeasible path

```
int global;  
int func() {  
    int undef_var;  
    ...  
    if (undef_var) {  
        lock();  
    }  
    global++;  
    if (undef_var) {  
        unlock();  
    }  
}
```



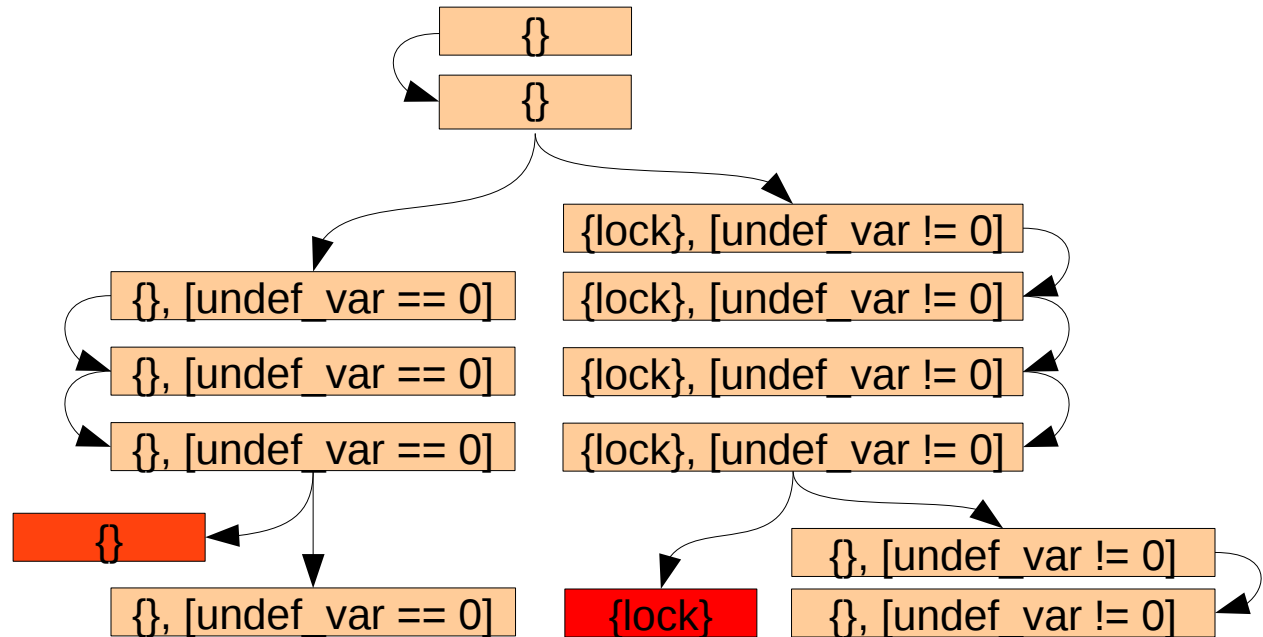
CEGAR

Counter Example Guided Abstraction Refinement



Example of infeasible path

```
int global;  
int func() {  
    int undef_var;  
    ...  
    if (undef_var) {  
        lock();  
    }  
    global++;  
    if (undef_var) {  
        unlock();  
    }  
}
```



Based on open-source projects

- Linux Driver Verification (LDV)
<http://linuxtesting.org/project/ldv/>
- CPAchecker
<http://cpachecker.sosy-lab.org>



Thank you

Questions?