
Comparing business models using graph-edit distance

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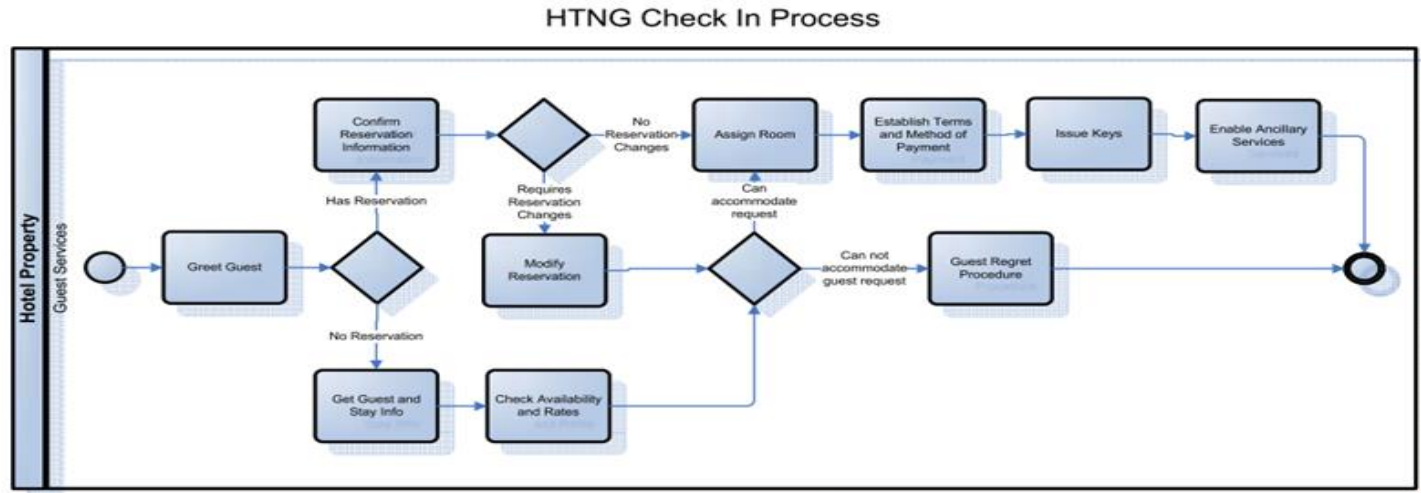
Agenda

- Business process models (BP)
- Changes in BPs
- BPMN
- BPMN model as a graph
- Graph-edit distance
- Our approach for comparing two BP models

- 40 mins
- Q&A on the spot

BP models

Business process modeling (BPM) in systems engineering is the activity of representing processes of an enterprise, so that the current process may be analyzed or improved.

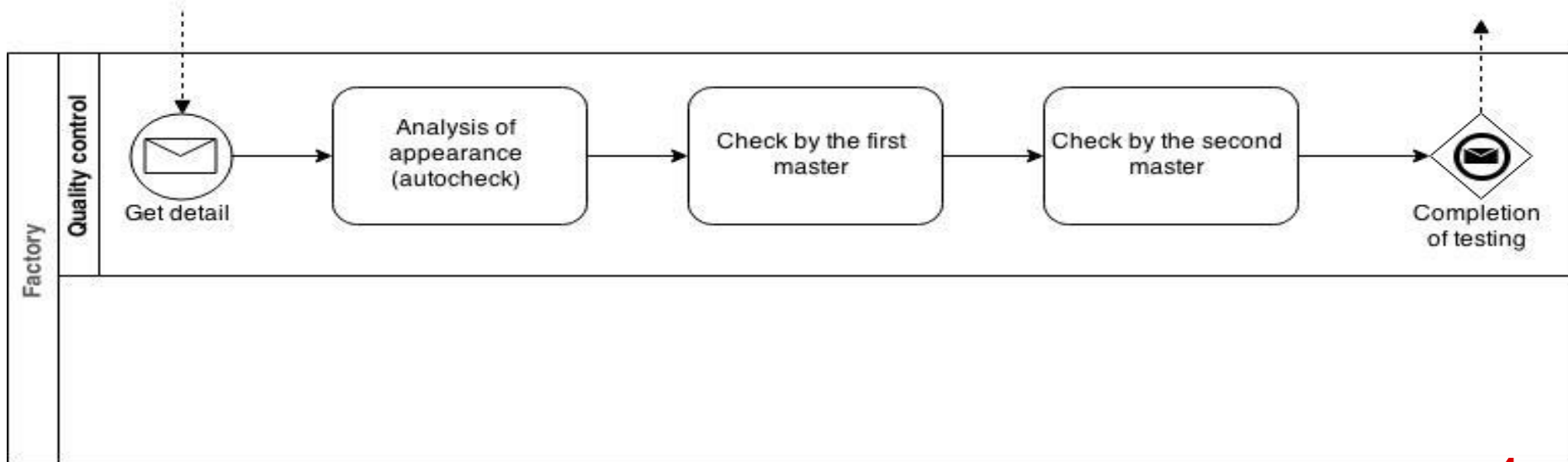


Changes in business processes

- new amendments to the legislation
- changes in revenue sources
- increase production efficiency
- economic progress
- technological progress

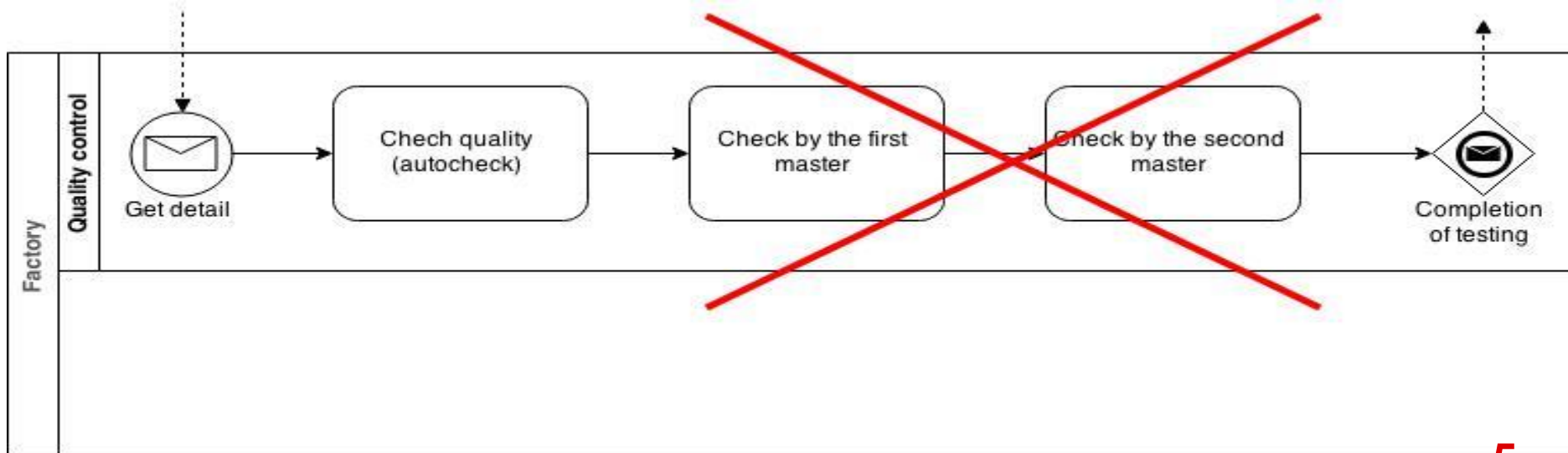
Example 1-1

Double quality control by two masters



Example 1-2

Purchase of equipment for testing quality of details



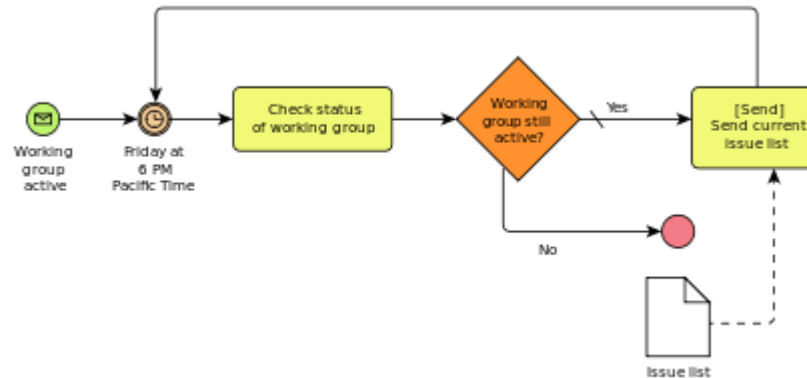
Business process modeling tools

- BPMN - Business Process Model and Notation
- CogNIAM - Cognition enhanced Natural language Information Analysis Method
- xBML - Extended Business Modeling Language
- EPC - Event-driven process chain
- IDEF0 - ICAM DEFinition
- UML - Unified Modeling Language



BPMN

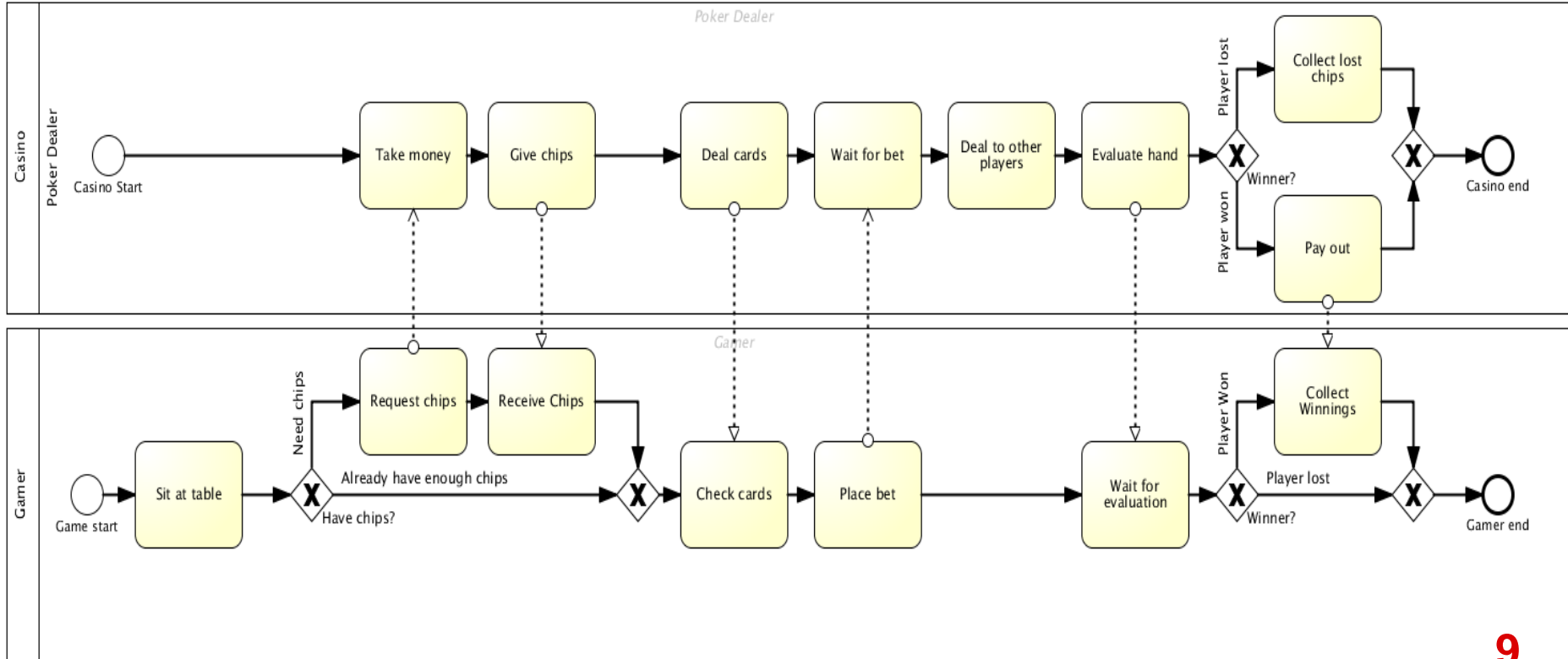
Business Process Model Notation (BPMN) is a graphical representation for specifying business processes in a business process model.



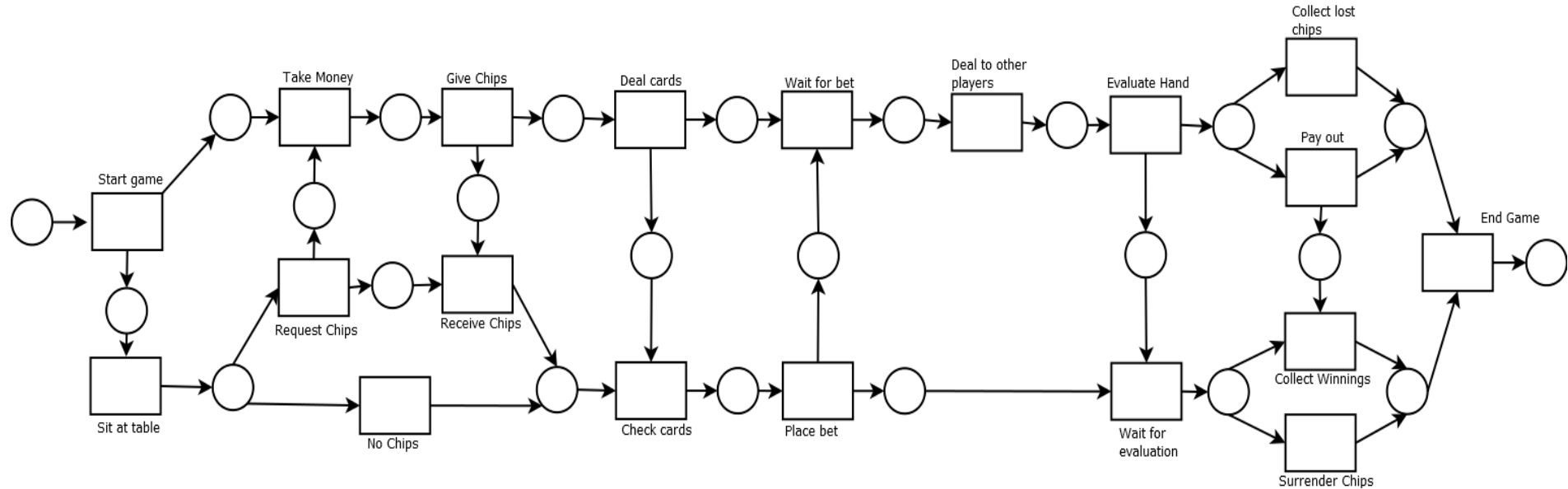
BPMN features

- Must be acceptable and usable by the business community for general process modeling
- Must be able to generate executable processes from a BPMN Model (through a combination of graphical elements and supporting information)
- BPMN is intended to be Methodology

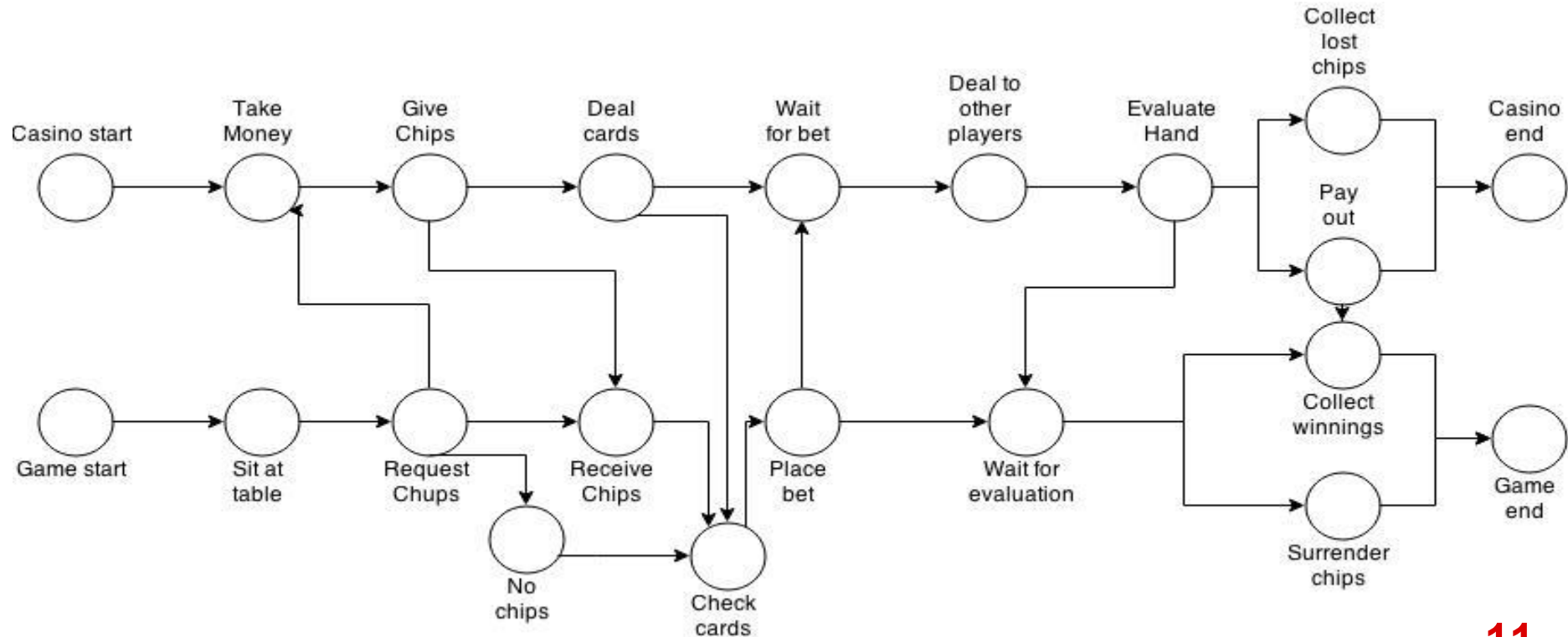
BPMN model



BPMN model as a net

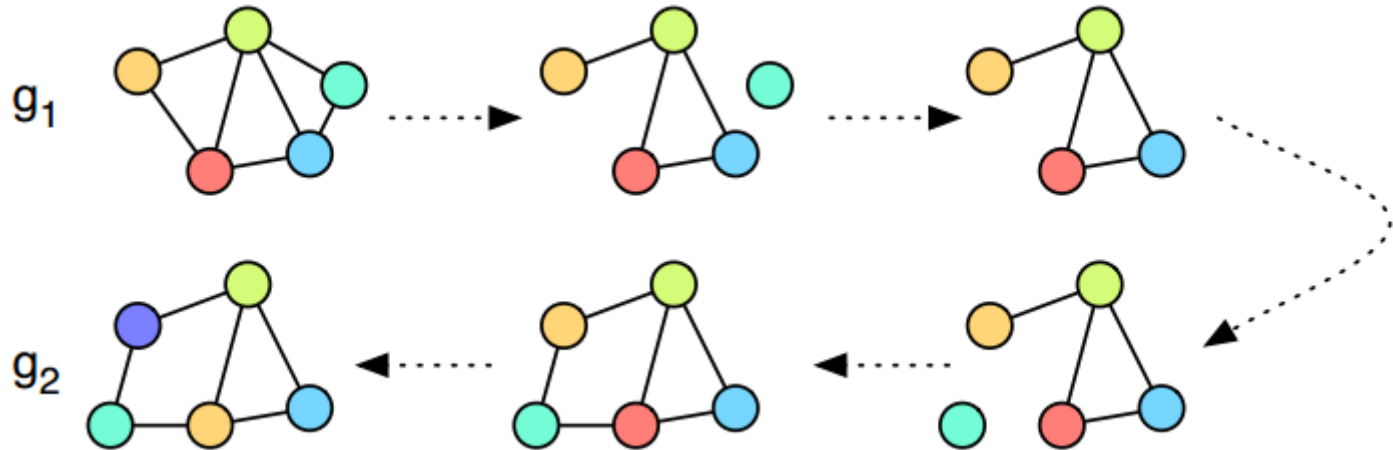


BPMN model as a graph



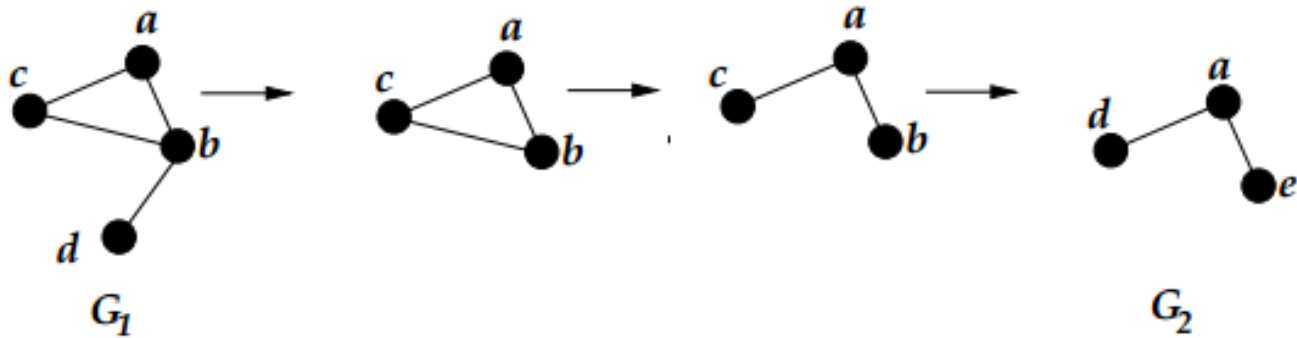
Comparing graphs

Calculate graph edit distance using A* algorithm



Graph edit distance

The graph edit distance d of two graphs g_1 and g_2 is the minimum costs equence of edit operations (next slide) that transform g_1 in to g_2 .



Graph edit operations

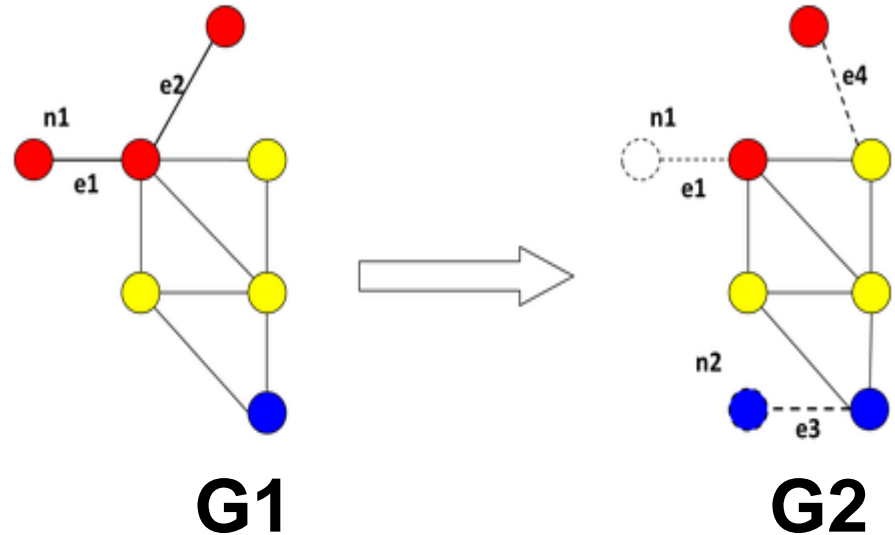
A graph edit operation is an edit operation to transform one graph to another, including 6 types:

- Insert an isolated vertex into the graph
- Delete an isolated vertex from the graph
- Change the label of a vertex
- Insert an edge between two disconnected vertices
- Delete an edge from the graph
- Change the label of an edge

Example 2-1

Cost of edit operations:

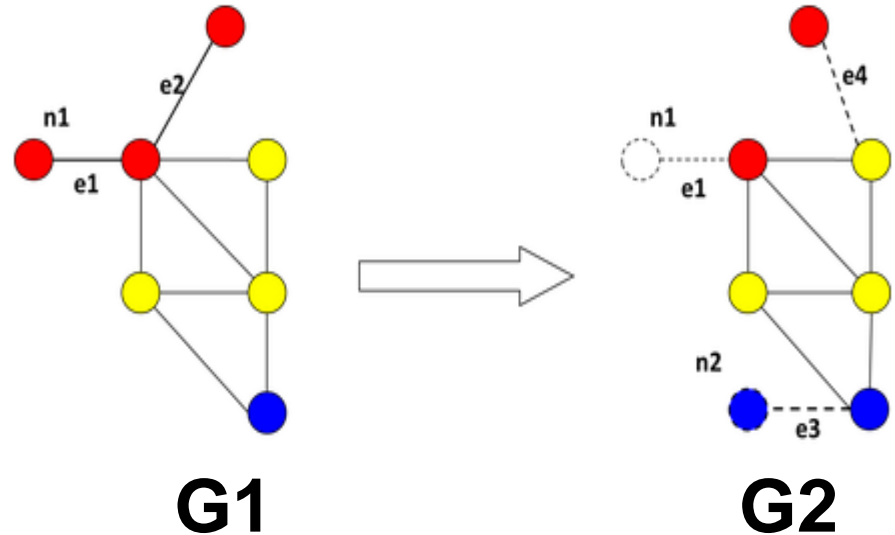
- deletion/insertion of edges/nodes – 1
- substitution of nodes/edges – 0.



Example 2-2

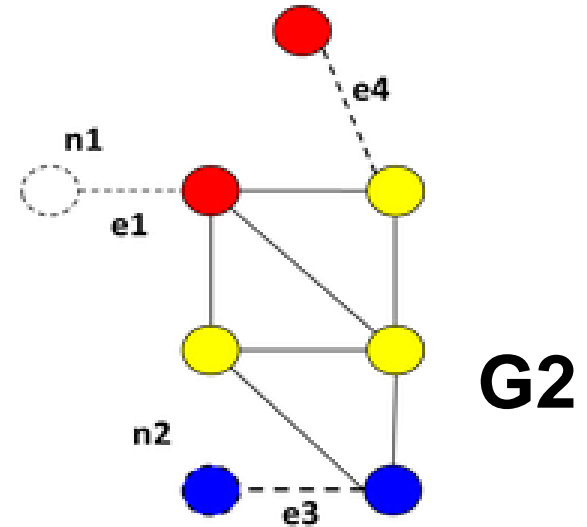
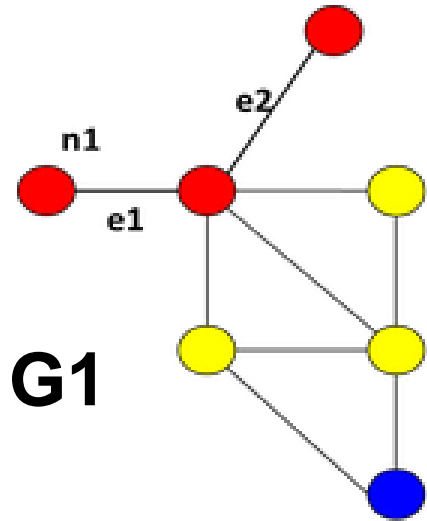
Transforming:

- delete edge **e1**
- delete node **n1**
- delete edge **e2**
- insert edge **e4**
- insert node **n2**
- insert edge **e3**



Example 2-3

$$\text{GED}(\mathbf{G1}, \mathbf{G2}) = 6$$



Graph edit distance

Two main drawbacks:

- a high computational complexity: the problem of computing graph edit distance is NP-hard in general
- the difficulty related to defining cost functions

Possible solutions:

- the most known method for computing the exact value of graph edit distance is based on A* search algorithm
- experimental evaluations

Graph edit distance by A* algorithm

Algorithm 1. Computation of graph edit distance by A* algorithm

Input: Non-empty graphs $g_1 = (V_1, E_1, \mu_1, \nu_1)$ and $g_2 = (V_2, E_2, \mu_2, \nu_2)$,
where $V_1 = \{u_1, \dots, u_{|V_1|}\}$ and $V_2 = \{v_1, \dots, v_{|V_2|}\}$

Output: A minimum-cost edit path from g_1 to g_2
e.g. $p_{min} = \{u_1 \rightarrow v_3, u_2 \rightarrow \varepsilon, \dots, \varepsilon \rightarrow v_6\}$

- 1: Initialize OPEN to the empty set
 - 2: For each vertex $w \in V_2$, insert the substitution $\{u_1 \rightarrow w\}$ into OPEN
 - 3: Insert the deletion $\{u_1 \rightarrow \varepsilon\}$ into OPEN
 - 4: **loop**
 - 5: Remove $p_{min} = \arg \min_{p \in \text{OPEN}} \{g(p) + h(p)\}$ from OPEN
 - 6: **if** p_{min} is a complete edit path **then**
 - 7: Return p_{min} as the solution
 - 8: **else**
 - 9: Let $p_{min} = \{u_1 \rightarrow v_{i_1}, \dots, u_k \rightarrow v_{i_k}\}$
 - 10: **if** $k < |V_1|$ **then**
 - 11: For each $w \in V_2 \setminus \{v_{i_1}, \dots, v_{i_k}\}$, insert $p_{min} \cup \{u_{k+1} \rightarrow w\}$ into OPEN
 - 12: Insert $p_{min} \cup \{u_{k+1} \rightarrow \varepsilon\}$ into OPEN
 - 13: **else**
 - 14: Insert $p_{min} \cup \bigcup_{w \in V_2 \setminus \{v_{i_1}, \dots, v_{i_k}\}} \{\varepsilon \rightarrow w\}$ into OPEN
 - 15: **end if**
 - 16: **end if**
 - 17: **end loop**
-

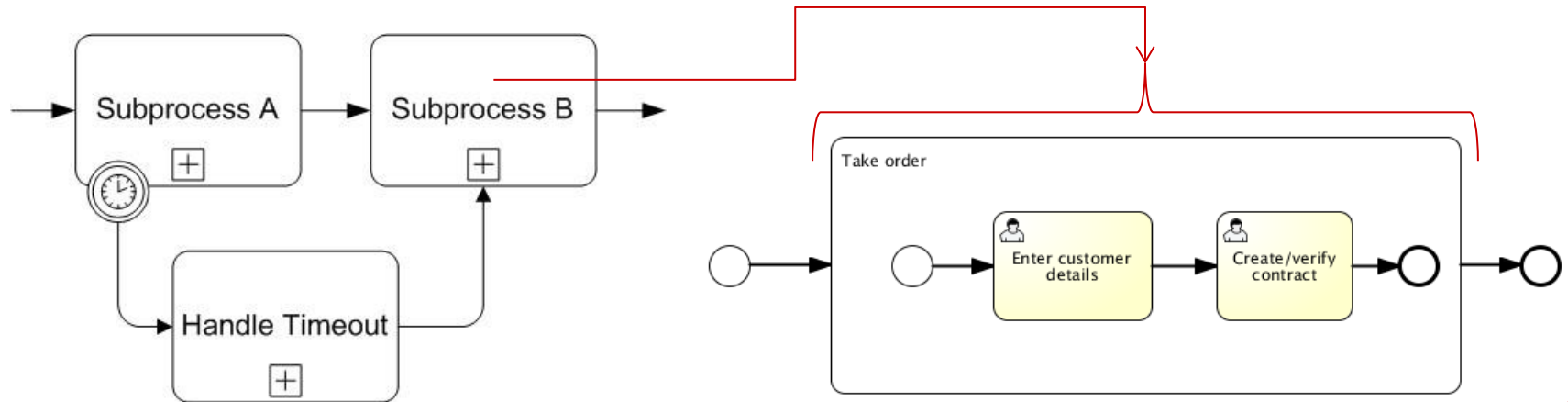
Comparing BPMN models

Differences from the comparison graphs:

- each element has a label
- each element has a special type (event, activity, gateway, connection)
- subprocess like subgraphs

Comparing BPMN models

A subprocess element should be represented as a subgraph which means that nested elements should be considered as separate elements.



Comparing BPMN models

Comparing two elements:

- string-edit distance for labels
- comparing types

Comparing two models:

- string-edit distance + comparing types
- A* algorithm

Comparing two BPMN models:

- string-edit distance + *preprocessing*
- A* algorithm + high-level comparison

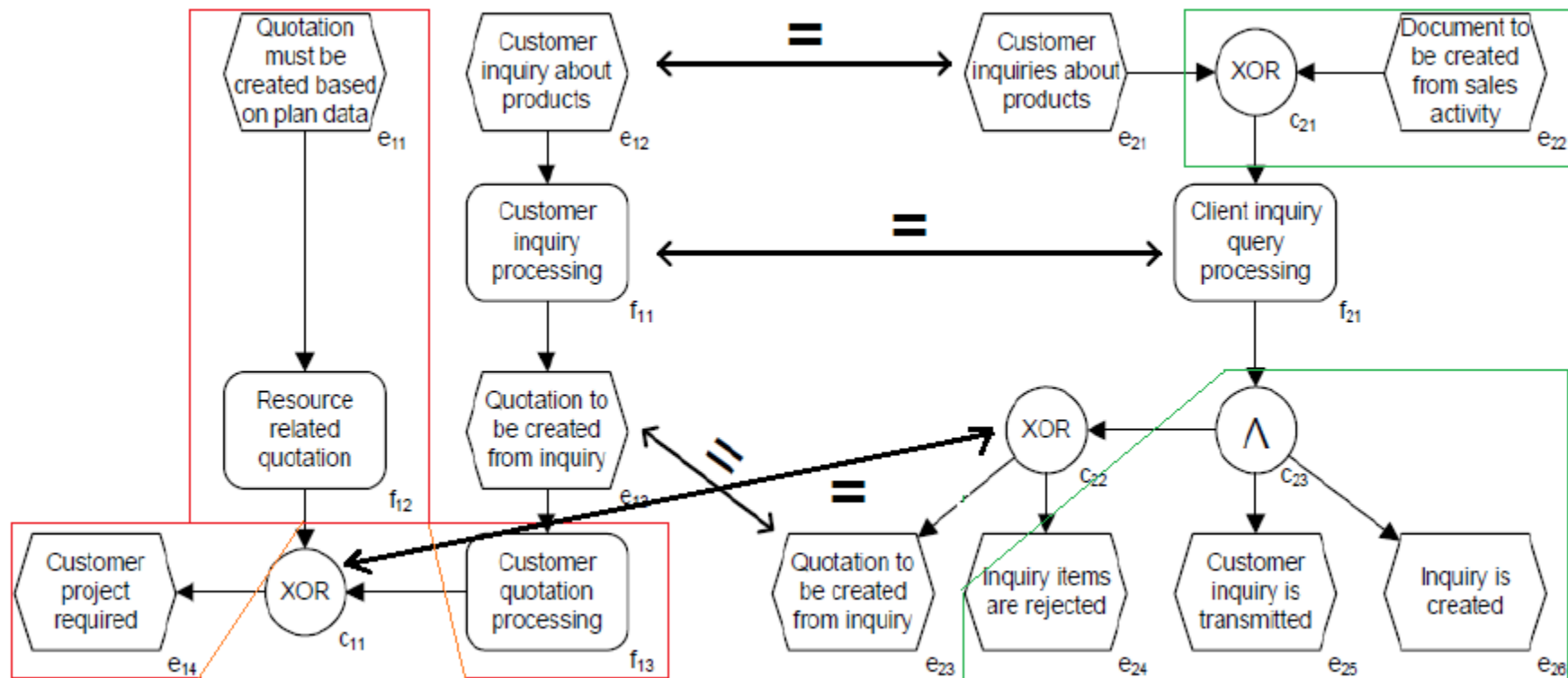
String edit distance distance

The Levenshtein distance is a string metric for measuring the difference between two sequences.

Operations: insertion, deletion or substitution.

1. kitten → sitten (substitution of "s" for "k")
2. sitten → sittin (substitution of "i" for "e")
3. sittin → sitting (insertion of "g" at the end).

Example 3-1

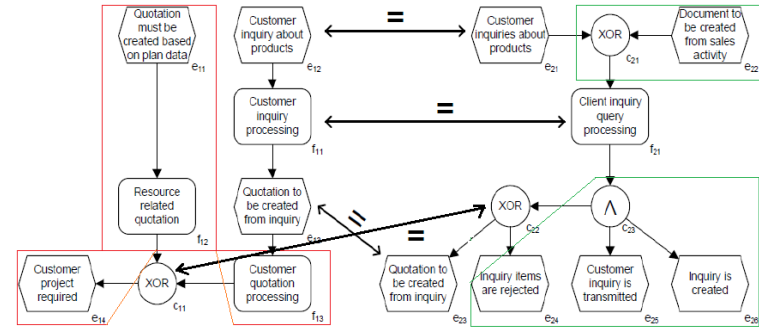


Example 3-2

Graph edit distance = 14

Actions:

- delete event 1
- insert activity 5
- change label for event 3
- ...



**Thanks a lot
for your attention!**
